



Towards an Error Control Scheme for a Publish/Subscribe Network

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Outline

- Motivation
- Background
- Operation
- Simulator setup
- Performance
- Conclusions

Motivation

- ICN: Information Centric Networking
 - ◆ Focuses on information rather than on endpoints
 - ◆ Most importantly, it supports native multicast
- Reliable multicast transport over ICN
 - ◆ Some multicast applications do not require reliability
 - Example: Multimedia streaming
 - ◆ But others really depend on it
 - Example: Software updates, sensor readings
 - ◆ Lots of work on reliable IP multicast, e.g. PGM
 - ◆ We adapt these ideas to the ICN case
 - Specifically, over the PSI architecture

Background

- Publish Subscribe Internet (PSI) architecture
 - ◆ Publishers advertise available data
 - ◆ Subscribers express interest in data
 - ◆ The Topology Manager creates paths between them
- Stateless forwarding in PSI
 - ◆ Paths are encoded as source routes
 - ◆ A Bloom filter includes the corresponding links
- Relay points used for forwarding scalability
 - ◆ Bloom filters cannot cover large trees
 - ◆ So we break them into connected subtrees
 - ◆ Relay points used to switch Bloom filters

Operation

- NACK based protocol
 - ◆ Only missing packets are NACKed
 - ◆ Relay points serve as NACK aggregation points
- Reverse Bloom filters used for the NACKs
 - ◆ Simply use the reverse links of the subtree
- Phase 1: Setup
 - ◆ Calculation of necessary Bloom filters
 - Both forward and reverse, for each subtree
 - ◆ Initial message from publisher
 - Propagates downstream
 - Relay points store Bloom filters

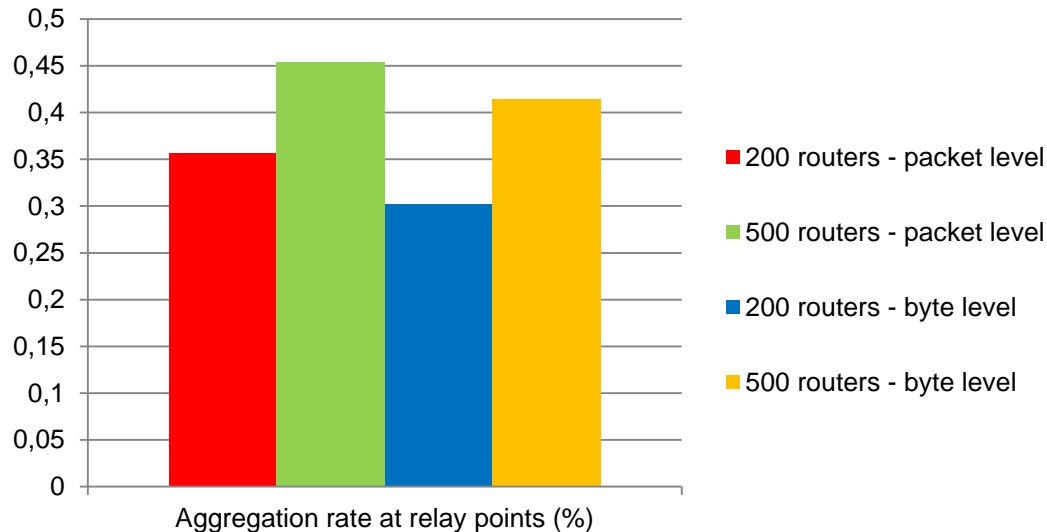
Operation

- Phase 2: Initial content distribution
 - ◆ Publisher sends entire content
 - ◆ Receivers send NACKs for missing packets
- Phase 3: Recovery
 - ◆ Repeats Phases 1 and 2
 - ◆ Some subscribers leave at the end of each cycle
 - ◆ Eventually no subscribers left
- NACK aggregation
 - ◆ Relay points receive NACKs
 - ◆ Wait to get more NACKs for a little while
 - ◆ Then, received NACKs are merged and propagated

Simulator setup

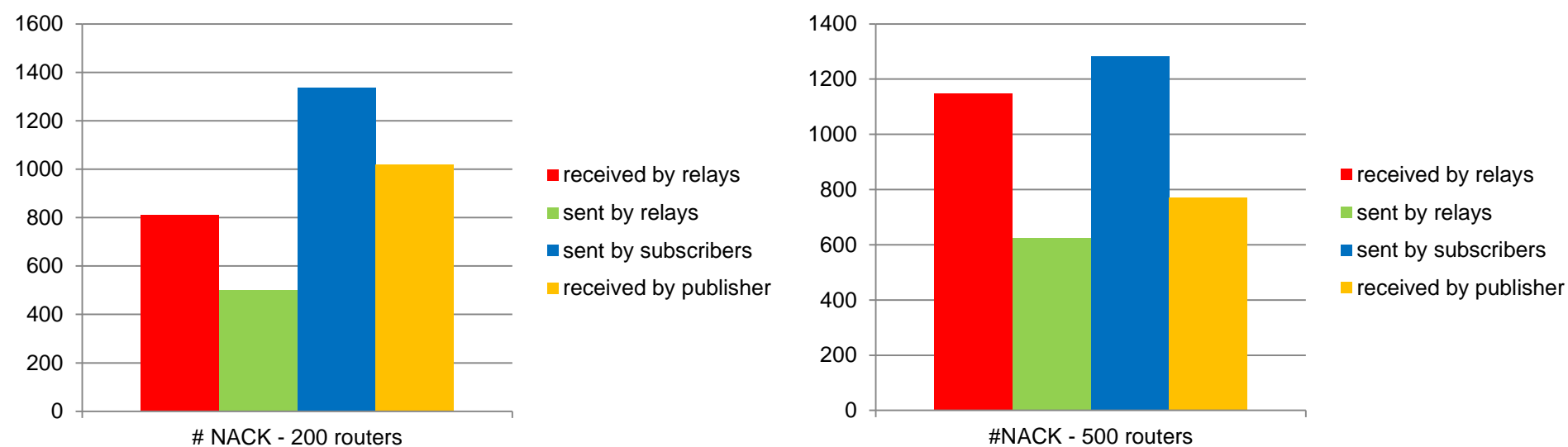
- NS-3 based simulations
 - ◆ Entire PSI architecture implemented
 - ◆ Relaying and transport protocol added
 - ◆ Single publisher transmits 20 MB
 - ◆ Scale free topologies with 200 and 500 routers
 - ◆ 50 and 100 subscribers randomly attached to routers
 - ◆ 3% of the packets are reported lost in first round
 - Random loss model used
 - Roughly 600 packets need to be retransmitted
 - Manual setting of link loss to achieve target rate

Performance: aggregation rate



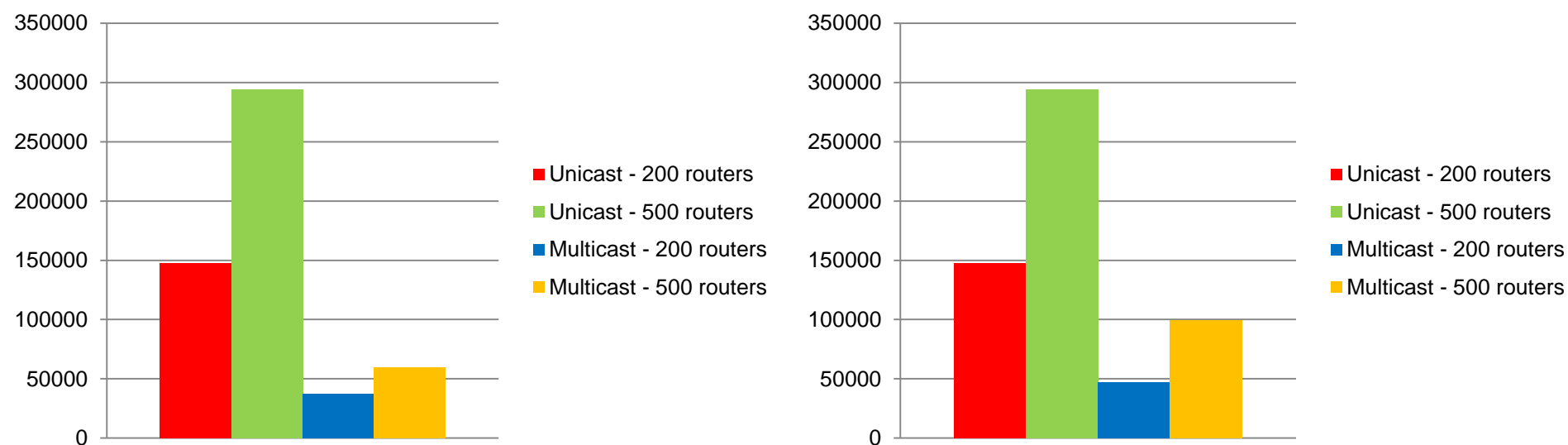
- Aggregation rate of NACKs
 - ◆ Fraction of NACKs *not* reaching the publisher
 - ◆ Grows with larger topology
 - More opportunities for aggregation
 - ◆ Slightly lower in bytes

Performance: NACKs handled



- Number of NACKs handled by each entity
 - ◆ Similar number of original NACKs
 - Both topologies start with the same loss rate
 - ◆ More NACKs handled by relays with larger topology
 - ◆ But, more aggregation overall with larger topology

Performance: Uni vs. Multicast



- Unicast vs. Multicast recovery
 - ◆ NACKs and retransmissions
 - Calculated over all network elements
 - ◆ Left: Number of NACK transmissions
 - ◆ Right: Number of packet retransmissions

Conclusions

- Reliable multicast transport in PSI
 - ◆ Takes advantage of relay points
 - ◆ Very good feedback aggregation
 - ◆ Beats unicast by a wide margin
- Ongoing work
 - ◆ Detailed comparison with PGM
 - ◆ PGM exploits unidirectional multicast
 - Relay point multicasts NACK confirm
 - Good for “correlated” losses
 - ◆ We exploit reverse Bloom filters
 - Relay point just aggregates NACKs
 - Good for “uncorrelated” losses