



# *Reducing Forwarding State in Content-Centric Networks with Semi-Stateless Forwarding*

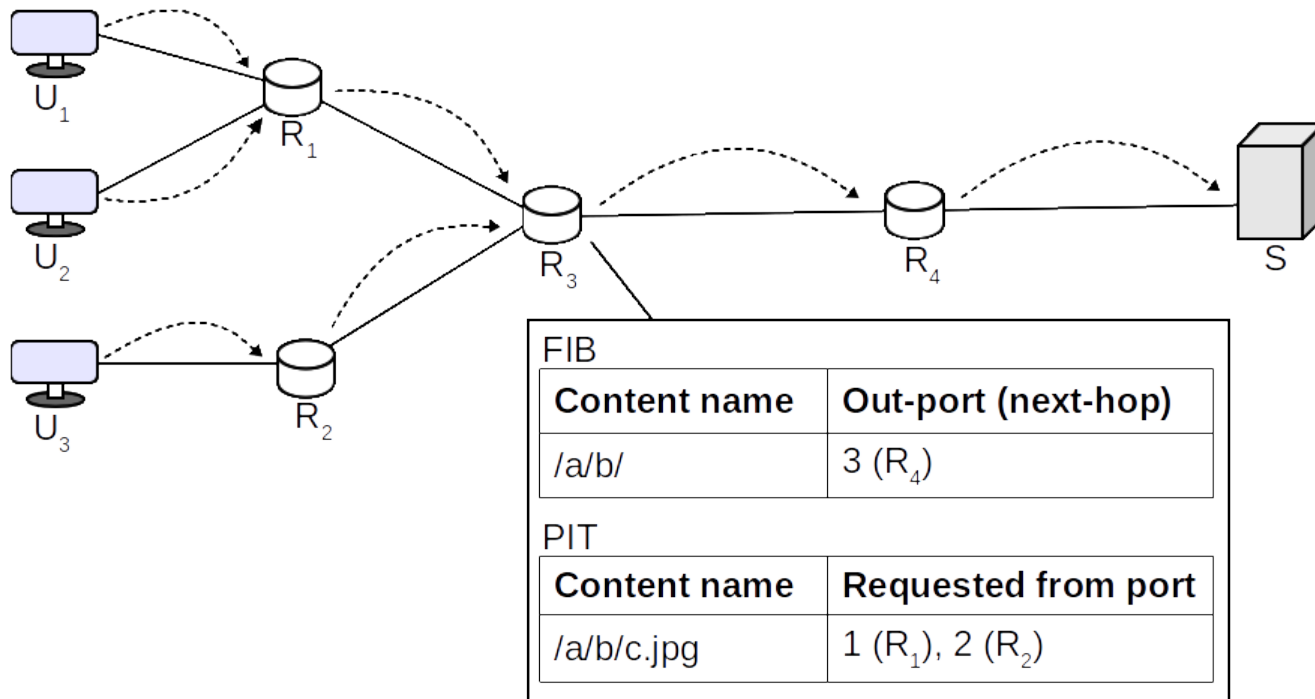
Christos Tsilopoulos, George Xylomenos, Yannis Thomas

Mobile Multimedia Laboratory,  
Athens University of Economics and Business

# Outline

- The problem: Does the PIT scale in CCN?
  - As is, very unlikely!
- A semi-stateless forwarding scheme
  - Interest tracking
  - Bloom filter-based Data forwarding
- Evaluation
  - Unicast & Multicast
- Conclusions

# A CCN Primer



- Hop-by-hop Interest forwarding (FIB)
- Pending Interests Table (PIT)
- Stateful forwarding advantages

# PIT Scalability

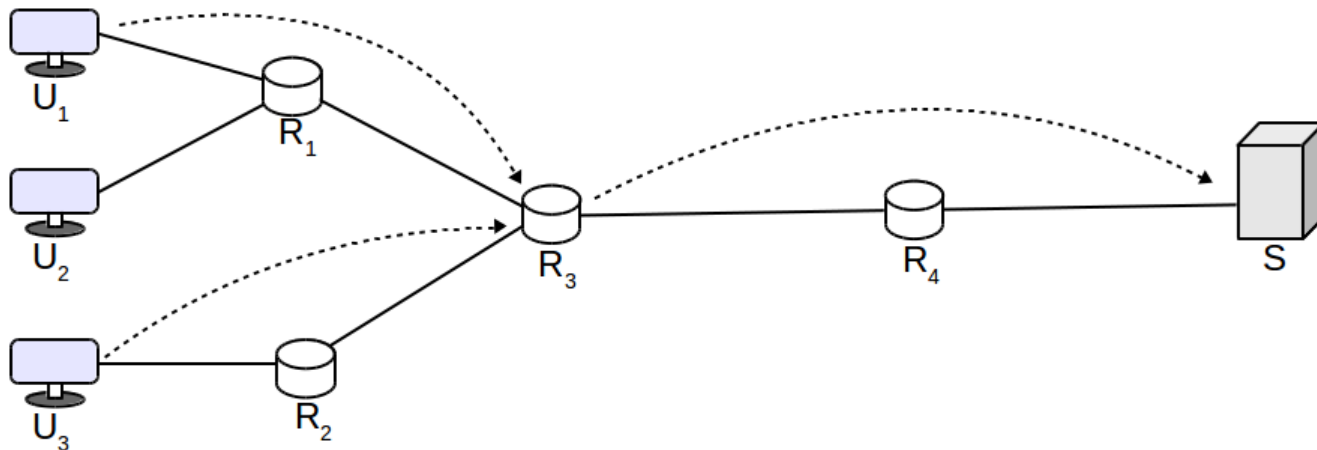
- # entries:  $\text{bandwidth} \times \text{RTT} / \text{packet}$ 
  - Example: 40 Gbps, 80 ms, 1 Kbyte
  - 400K PIT entries, longer than IP addresses
  - Even more in live streaming and pub/sub
  - Realistic estimates indicate millions of entries
- Large memory requirements for PIT
  - Does not fit in on-chip memory
  - Bad performance with main memory

# Reducing PIT size

- DiPIT: combine entries with Bloom filters
  - Drops per Interest information
- ENPT: uses a trie to encode Interests
  - Lookup time depends on name size
- CONET: Interests accumulate source routes
  - Drops all benefits of stateful forwarding
  - No multicast or adaptive forwarding
  - No dropping of unwanted Data

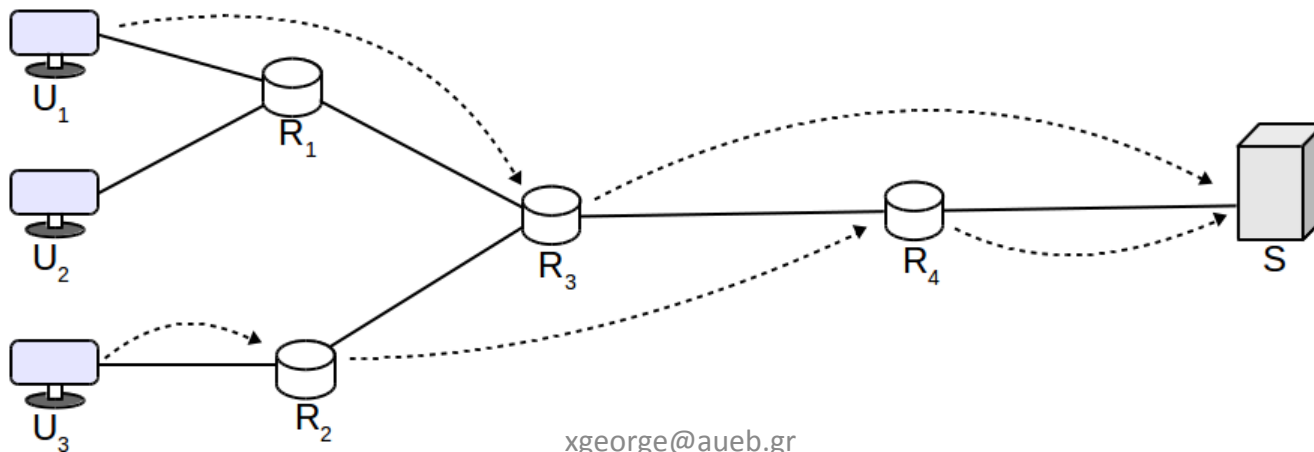
# Semi-Stateless Forwarding

- Track Interests at *some* on-path routers
  - On average, every  $d$  hops
  - For  $N$ -hop paths, Interest tracked at  $N/d$  routers
- Stateless forwarding between stateful routers
  - Using in-packet Bloom filters

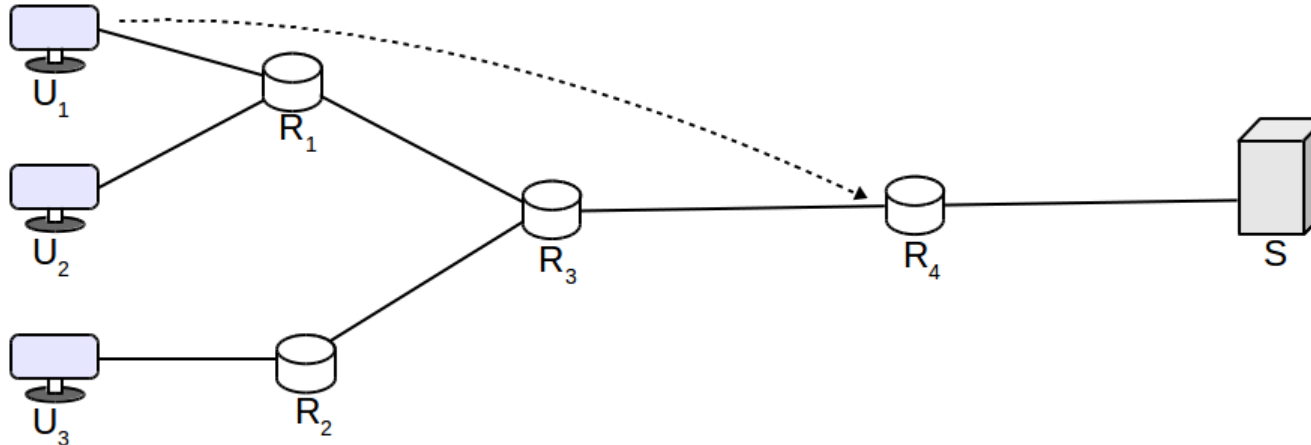


# Interest Tracking: Requirements

- Spread state across routers
  - Avoid bottleneck points
  - Routers should track  $1/d$  of forwarded Interests
- Efficient multicast rendezvous
  - Aggregate Interests for the Same Data



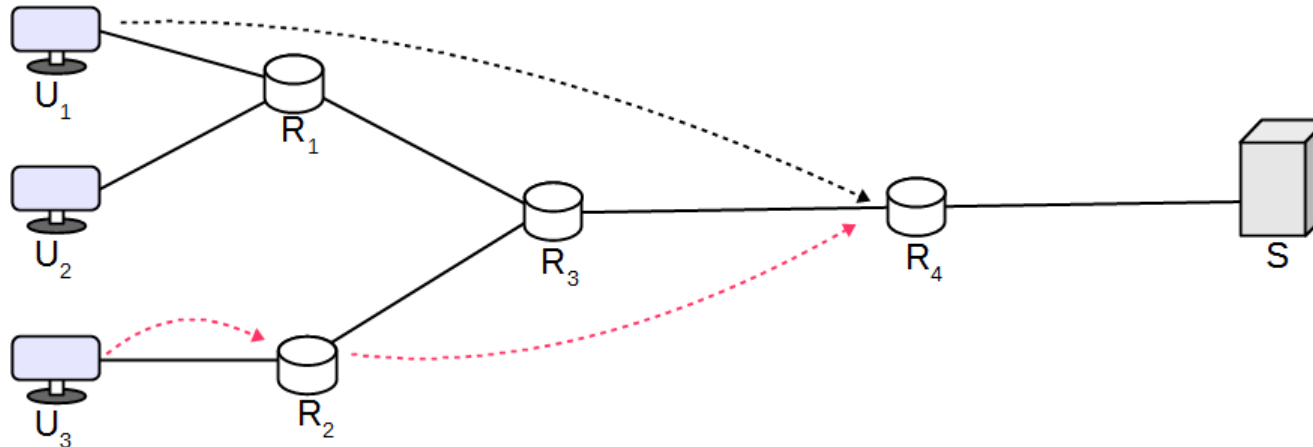
# 1. Probabilistic Tracking



- Router tracks Interest with probability  $1/d$ 
  - Interest aggregation: check if entry exists
  - Upper bound on stateless hops
    - Avoid large stateless parts
    - Uses hop counter in Interests

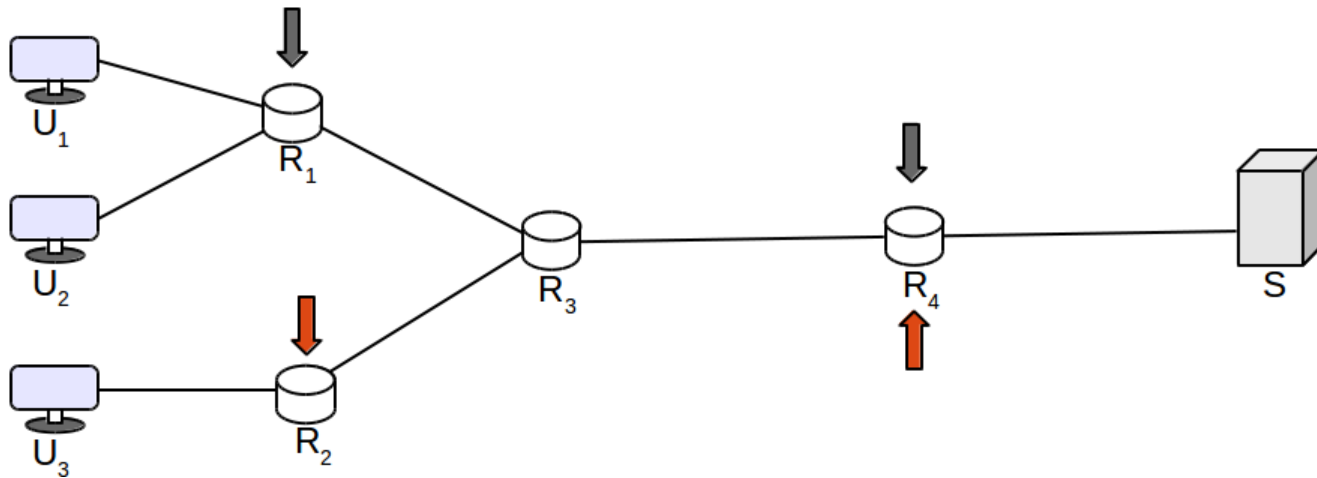


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## 2. Hash-based Tracking

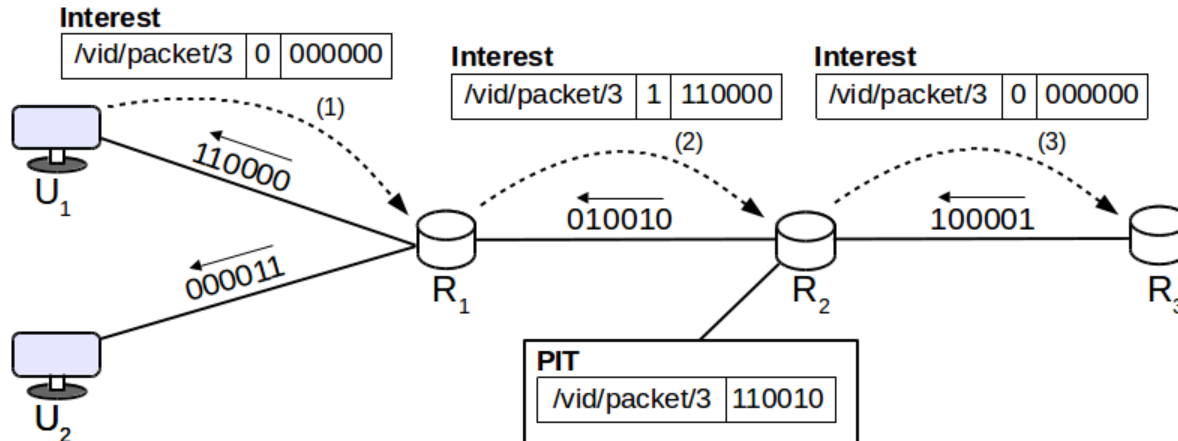


- Router tracks Interest of  $h \bmod d == 0$ 
  - $h = \text{hash}(\text{content\_name} + \text{router\_suffix})$
  - Deterministic selection of storage points
  - Also needs upper bound on stateless hops

# 3. Hop Counter-based Tracking

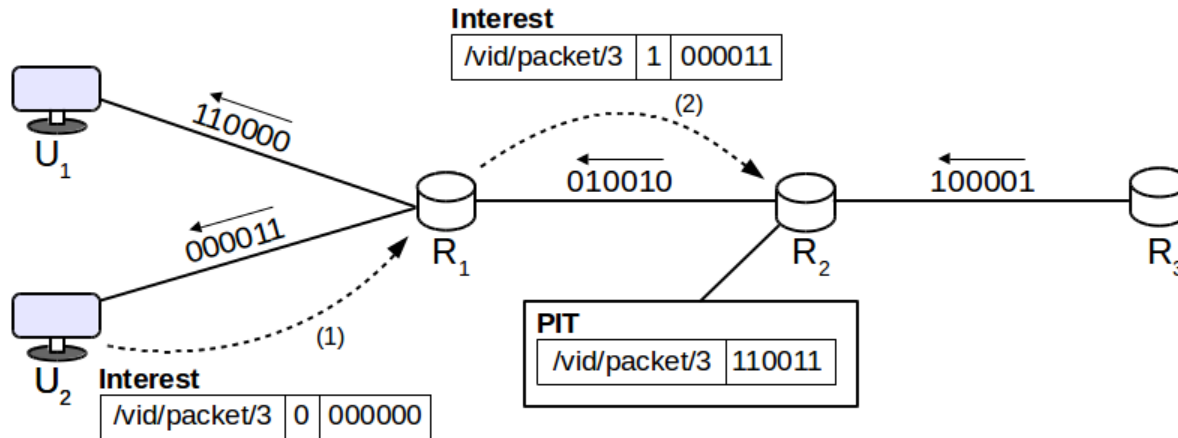
- Both previous policies require hop counters
- Why not use *only* the hop counters?
  - Much simpler to implement
  - Stateless paths never get long
- Select initial value randomly in  $[0, d-1]$ 
  - Spread state across routers
- Check PIT for pre-existing entry
  - Aggregate entries as soon as possible

# Data Forwarding: Bloom filters



- Track reverse path information in routers
  - Each link has a random ID (LID)
  - Add (OR) reverse LID at each hop
  - Store the Bloom filter (BF) in the PIT
  - Add (OR) the BFs to aggregate paths

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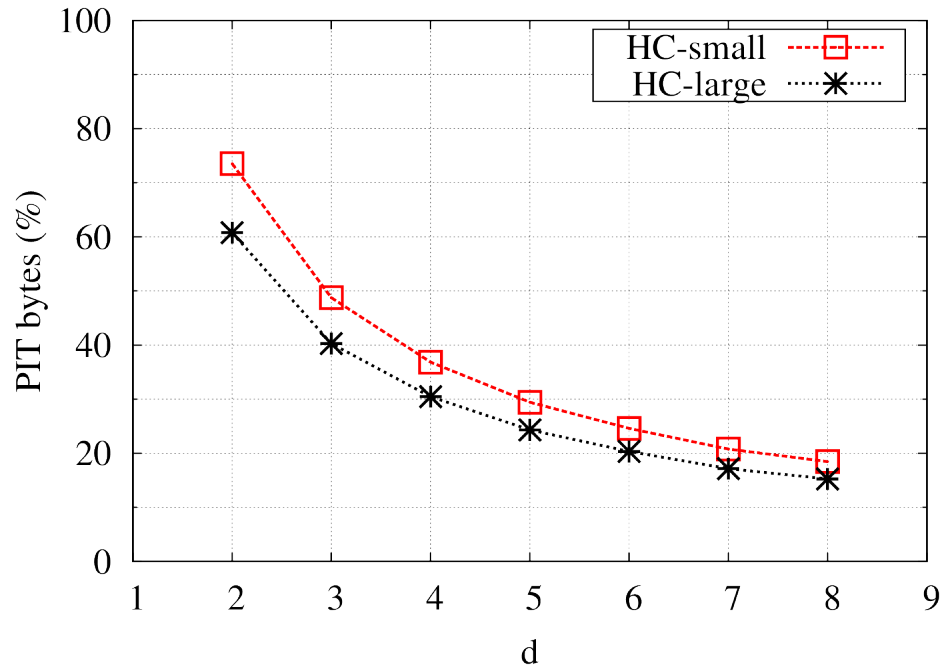
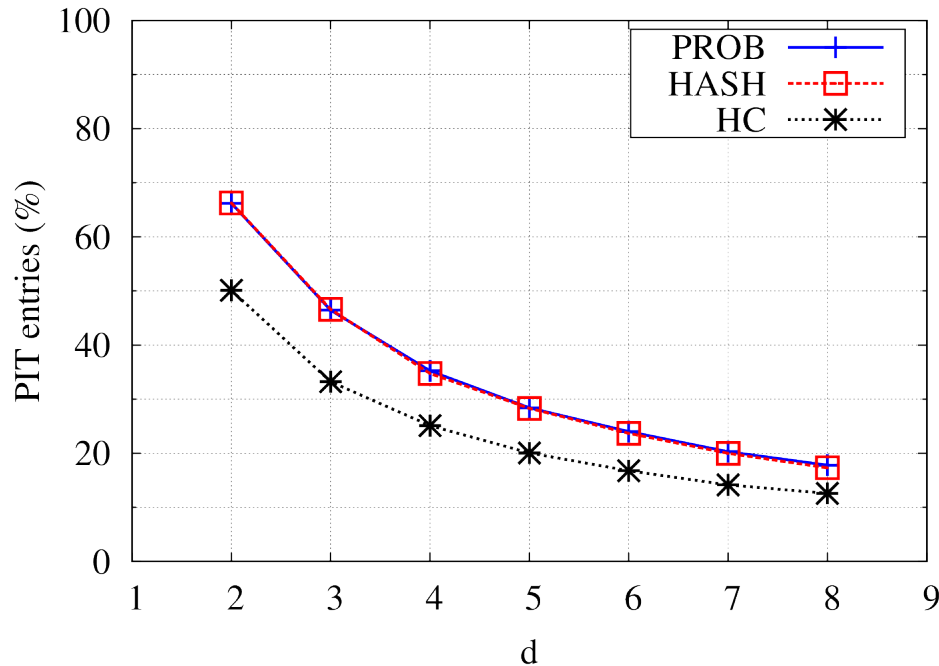
# Tradeoffs

- Larger Interest & Data packets
  - Need to carry BFs in headers
- Additional Interests with multicast
  - Aggregation beyond first common router
  - Does not influence Data transmissions
- Additional Data
  - False positives in BFs
- Fewer but larger PIT entries

# Evaluation

- Simulations using AS 224
  - Similar results with other ASes
  - 500 hosts distributed to access routers
  - 128 bit Bloom filters for forwarding
- Metrics
  - Reduction in PIT entries
  - Reduction in PIT size (bytes)
  - Additional Interests and Data

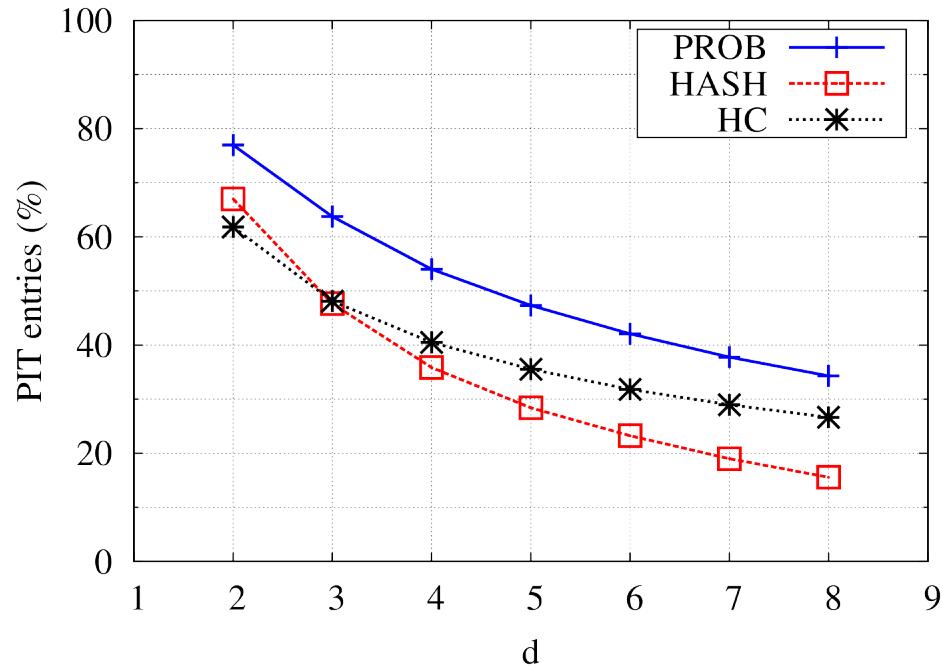
# Unicast



- Simple file transfer
  - 66% reduction in PIT entries (HC, d=3)
  - 54-61% reduction in PIT size (HC, d=3, small/large)

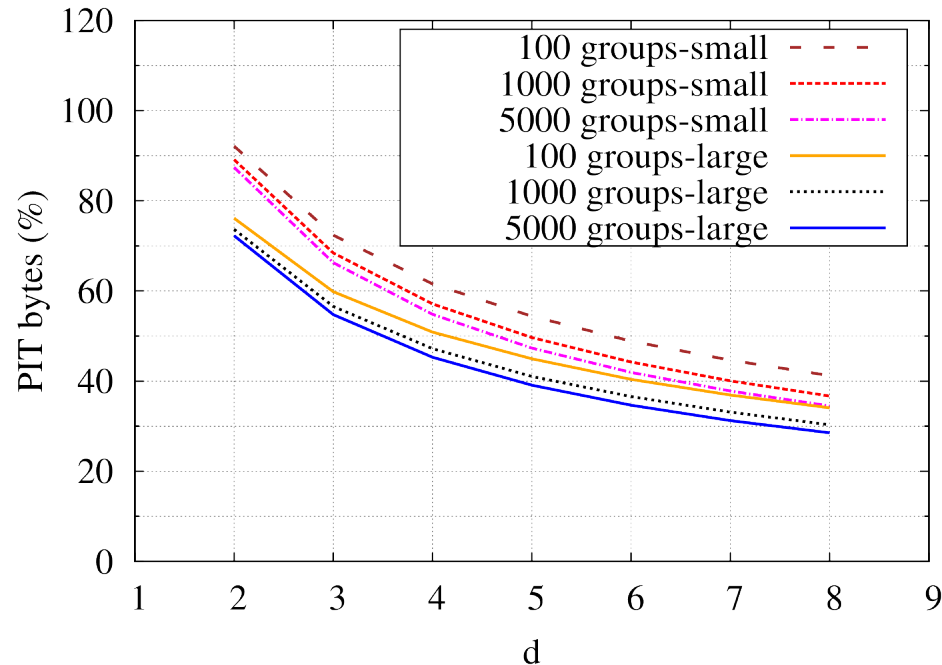


# Multicast: PIT entries



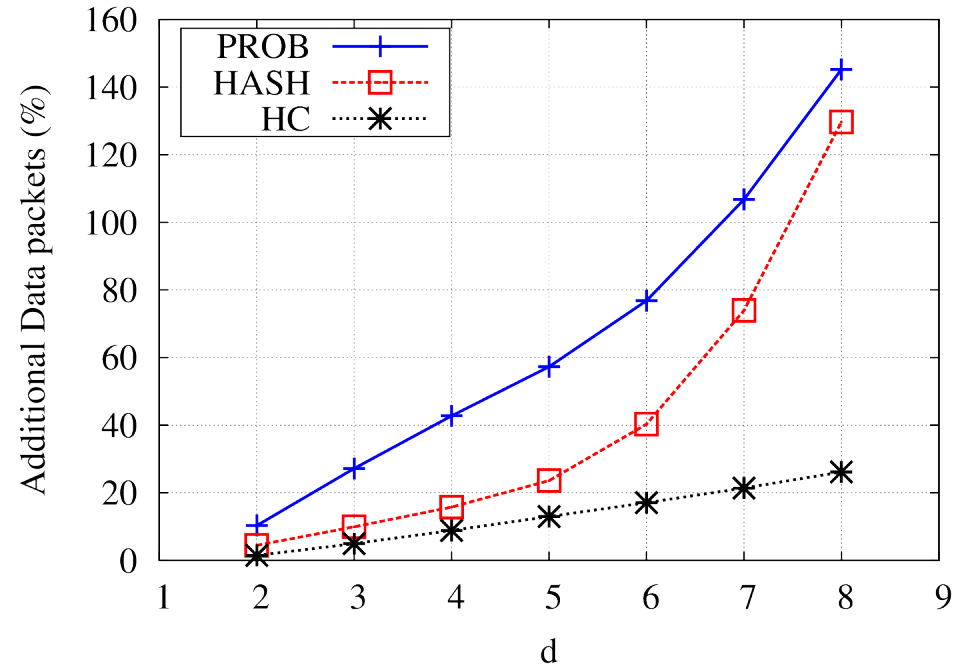
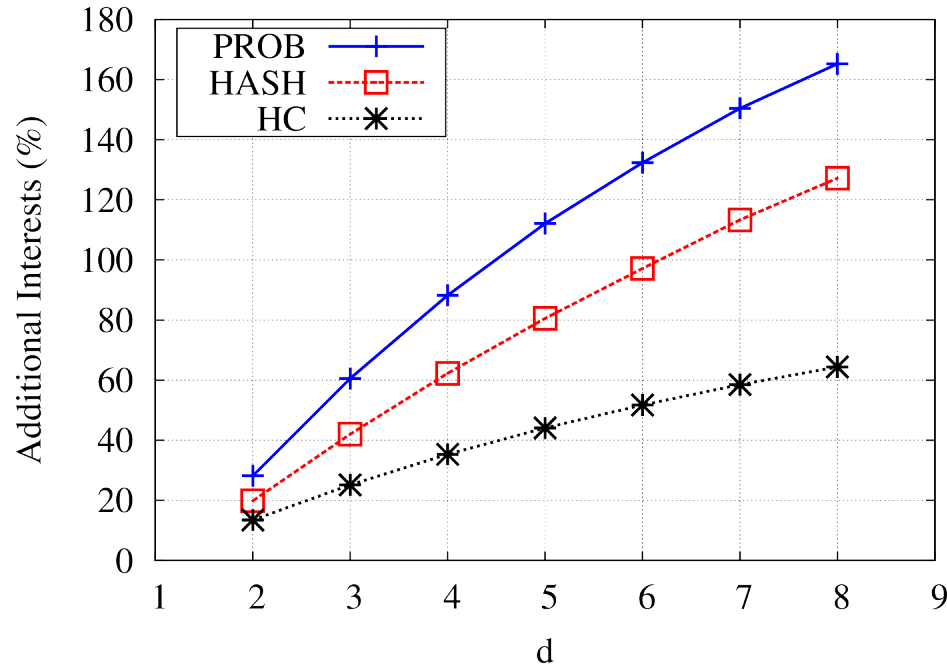
- Live media streaming
  - 1000 groups, Zipf distribution of sizes
  - 52% reduction in PIT entries (HC,  $d=3$ )

# Multicast: PIT Size



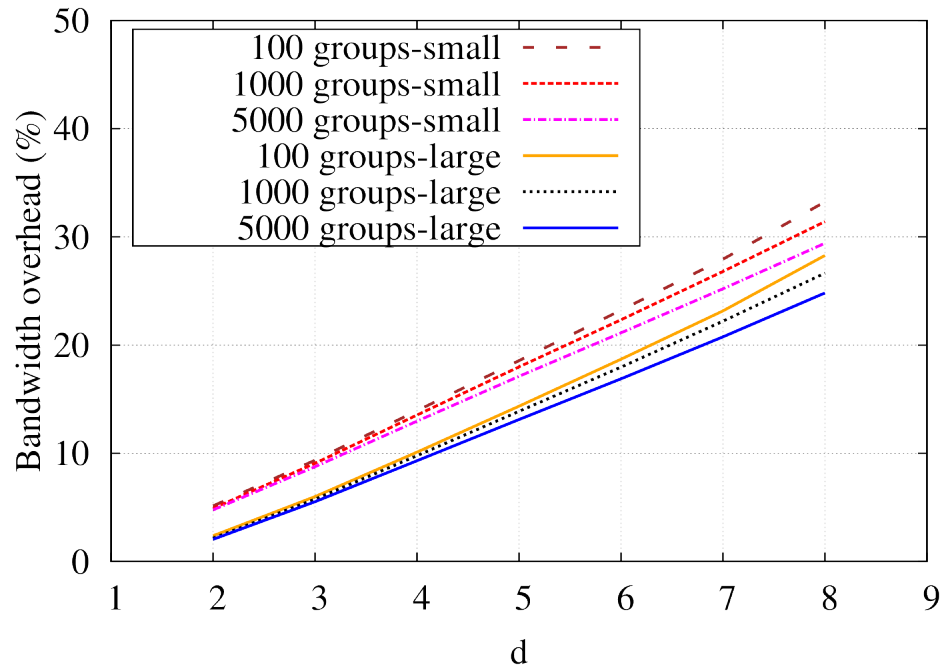
- Live media streaming
  - 34% smaller PIT size (5k groups, d=3, small names)
  - 45% smaller PIT size (5k groups, d=3, large names)

# Multicast: Overhead



- Live media streaming
  - 25% additional Interests (HC,  $d=3$ )
  - 5% additional Data (HC,  $d=3$ )

# Multicast: Total Overhead



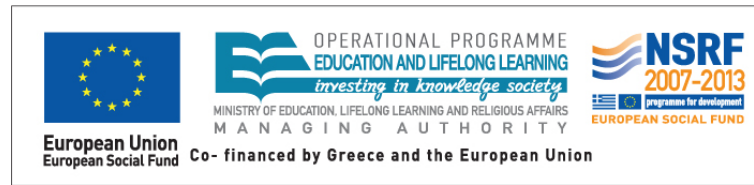
- Live media streaming
  - 9% more bandwidth (5K groups, d=3, 1.5K pkt)
  - 6% more bandwidth (5K groups, d=3, 8K pkt)

# Evaluation overview

- Significant PIT size reduction achieved
  - 54-61% compared to CCN for unicast
  - 34-45% compared to CCN for unicast
- Small bandwidth penalty
  - No penalty at all for unicast
  - 6-9% extra bandwidth on multicast
- The HC policy works best with  $d=3$  or 4
  - $d$  trades off PIT size against bandwidth overhead

# Conclusion

- Semi-stateful forwarding reduces PIT size
- Qualitative aspects of CCN are maintained
  - Multicast, dropping unwanted packets
- Only Interest & Data processing changes
  - LIDs are constructed autonomously
- Future work
  - Resort to semi-stateless only as PIT fills up
  - Vary  $d$  per router to minimize overhead



# Thank you

xgeorge@aueb.gr